

**cs70, fall 2001  
midterm 2 solutions  
professor wagner**

**Problem #1 (18 pts.) Short-answer questions**

(a)  $C(n,i)/2^n$

(b)  $0.3 \leq \Pr[E \cup F] \leq 0.5$ . 0.3 occurs if E is a subset of F, and 0.5 occurs if E and F are disjoint.

(c) 2002, since each such string is of the form  $[k]1[2000-k]0$  (where the quantity in brackets indicates the number of times to repeat) for some  $k$  in  $\{0, 1, \dots, 2001\}$ , and there are 2002 such  $k$ .

(d) All of them, since  $2z = n+1 = 1 \pmod n$ , so  $z$ -inverse =  $2 \pmod n$ . (Alternate explanation:  $\gcd((n+1)/2, n) = \gcd((n+1)/2, (n-1)/2) = \gcd(1, (n-1)/2) = 1$  by Euclidean algorithm, so  $\gcd(z, n) = 1$ , so  $z$  has an inverse mod  $n$ .)

**Problem #2 (12 pts.) Digit sums**

Consider an arbitrary natural number  $n$ .

Write  $n$  in decimal:  $n = A_k \cdot 10^k + \dots + A_1 \cdot 10 + A_0$

Then

$$\begin{aligned} f(n) &= A_k^3 + \dots + A_1^3 + A_0^3 && \text{[by defn of } f \text{]} \\ &= A_k + \dots + A_1 + A_0 \\ &= A_k \cdot 10^k + \dots + A_1 \cdot 10 + A_0 && \text{[since } 10^j \equiv 1 \pmod 3 \text{ for all } j \text{]} \\ &= n \pmod 3 \end{aligned}$$

$n$  was arbitrary, so this must hold for all natural numbers.

**Problem #3 (12 pts.) Computing with polynomials**

(a)  $O(d^2(\lg p)^2)$  There are  $O(d^2)$  cross-terms, and each requires  $O((\lg p)^2)$  work to do a modular multiplication and addition.

(b)  $O(d(\lg p)^2)$  Computing the sequence  $1, u, u^2, \dots, u^d \pmod p$  takes  $d$  multiplications, then multiplying by  $A_i \pmod p$  takes  $d+1$  multiplications, plus  $d$  more additions.

(c) 1. Let  $v = f(u) \pmod p$  using part (b).

2. Return  $v^2 \pmod p$

Takes  $O(d(\lg p)^2 + (\lg p)^2) = O(d(\lg p)^2)$  time.

Note that computing  $f(x)^2$  first (using part (c)) then evaluating at  $u$  gives a slower algorithm.

### Problem #4 (8 pts.) Mystical polynomials

(a) No.  $f(6) = 57$ , which is divisible by 3 and hence not prime. (But note that  $f(3) = 3$ , which is prime, so  $f(3)$  is not a counterexample to  $f$  being mystical.)

(b)

which means that 3 is a divisor of  $f(3)$ . **Consequently,  $f(3)$  cannot be prime**, which implies that  $f$  is not mystical.

The bold statement is wrong. The case  $f(3) = 3$  is compatible with 3 dividing  $f(3)$ , and yet 3 is prime. Hence the underlined statement does not follow, and in fact the polynomial in part (a) gives an example where the reasoning breaks down. Worse still, the "theorem" is wrong:  $f(x) = 3$  is a counterexample (it is mystical).

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