

# UC Berkeley : CS61C (Garcia & Lustig) : Midterm part 1 : 2014-10-10

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## Question 1: Running in circles (25 min, 18 pts)

A *nibble* is half of a byte (4 bits). You'd like to implement `LoadNibble` in MAL MIPS, a function that takes one `uint32_t` argument `N` and returns the  $N^{\text{th}}$  nibble of memory in the lowest 4 bits of the return register (the other 28 bits should be 0). Note: The  $N^{\text{th}}$  nibble immediately follows the  $N-1^{\text{th}}$  nibble without overlapping; see box . The MIPS instruction `srlv` ("shift right variable") might be useful here; it operates like the `shamt`-based right-shift, except that its 3<sup>rd</sup> *register* argument is the variable amount to shift by.

1/2

a) What fraction of all the nibbles of memory can you access? \_\_\_\_\_

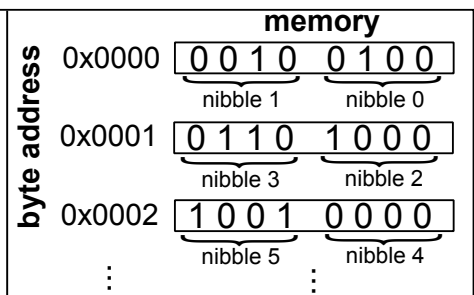
b) Implement `LoadNibble` by filling in the blanks:

```

LoadNibble:  srl  $a0  1
             lbu  $t0  $t0  # figure out which byte contains that nibble
             andi $a1  0(____)
             andi $a0  $a0  0x1
             sll  $a0  $a0  2 # we needed this!
gone1:      srlv  $v0  $a1  $a0
gone2:      andi  $v0  $v0  0xF
             jr   $ra
    
```

for N=2, `LoadNibble`  
returns `0b1000`

for N=5, `LoadNibble`  
returns `0b1001`



c) We want to rewrite `LoadNibble` to make use of a helper function `Helper` that will take two arguments. The first is an index `i` from 0-1 and the second is a byte `B`. `Helper` returns the `i`th nibble in `B` placed in the lowest 4 bits of the return value (the rest 0s).

E.g., `Helper(0, 0b01100100) → 0b0100` and `Helper(1, 0b01100100) → 0b0110`

We decide we don't need the two MIPS instructions labeled "gone1" and "gone2". What would you replace these instructions (and the `sll`) with to call `Helper` and implement `LoadNibble` successfully? Write the replacement below. Follow calling conventions and complete it in the fewest lines possible.

```

addiu $sp $sp -4
_____ # this line may not be necessary
sw $ra 0($sp)
_____ # this line may not be necessary
_____ # this line may not be necessary
jal _____
_____ # j works too, all other lines blank (since $ra = LoadNibble's caller)!
lw $ra 0($sp)
_____ # this line may not be necessary
addiu $sp $sp 4
_____ # this line may not be necessary
_____ # this line may not be necessary
    
```

## Question 2: *I can C clearly now, the rain is gone...* (25 min, 18 pts)

- A) Fill in the blank to complete this function that parses a string of *octal digits* (base 8) into a `uint64_t`. For example, calling `parse_octal("71")` should return the number 57. Do not use the comma operator, nested assignment, prefix/postfix operators, or function calls. You may assume that the given number "fits" into a `uint64_t`. (Hint: The backside of the MIPS green sheet may help.)

```
uint64_t parse_octal(char *s) {
    uint64_t r = 0;
    while(*s){
        r*8 + (*s - '0')
        r = _____;
        s++;
    }
    return r;
}
```

- B) We have the following data *packed tightly (no padding)* into the struct `data`, and some more code below:

```
struct {
    int16_t a;
    char b[2+(UNKNOWN_LENGTH*4)];
    int32_t c;
    int32_t d;
} data;
```

Fill in the blanks with an equivalent expression using only the pointer `s`, pointer arithmetic, casting, and the function `strlen()`. You may **NOT** use `UNKNOWN_LENGTH`. Assume `sizeof(char) = 1`.

```
/* ... Some code here that fills in data.b with the longest string possible ... */
char *s = data.b; /* s is a char, so it counts by 1 byte by default if in parens */
                s-1 /* or (s-2) */
*( (int16_t *) _____ ) = -1; // data.a = -1;
                (s+strlen(s)+1+4)
*( (int32_t *) _____ ) = -1; // data.d = -1;
```

- C) Here we have a *LR-tree*, defined as a node with two arrays of child pointers: two left children and two right children. Each node also contains a pointer to its parent node, a unique integer ID value, and a string name field. Root nodes will have a `NULL` parent pointer, and leaf nodes will have arrays of `NULL` children pointers.

```
struct lr_tree{
    char *name;
    uint64_t ID;
    struct lr_tree *left_children[2];
    struct lr_tree *right_children[2];
    struct lr_tree *parent;
};
```

Fill in the blanks to complete this function that frees a LR-tree if called with the root of the tree. You must free **ALL**

data associated with this LR-tree! You might not need all of the blanks, in which case use the most minimal number of blanks possible. Do not use the comma operator, nested assignment, or prefix/postfix operators.

```
void free_lr_tree (struct lr_tree *p) {
    p != NULL
    if ( _____ ){
        for(size_t x = 0; x < 2; x++) {
            free_lr_tree(p->left_children[x]);
            _____;
            free_lr_tree(p->right_children[x]);
            _____;
        }
        free(p->name);
        _____;
        free(p);
        _____;
    }
}
```