CS 61A Final Exam — May 15, 2009

Your name

login: cs61a–\_\_\_\_

This exam is worth 70 points, or about 23% of your total course grade. The exam contains 15 questions.

This booklet contains 14 numbered pages including the cover page. Put all answers on these pages, please; don't hand in stray pieces of paper. This is an open book exam.

When writing procedures, don't put in error checks. Assume that you will be given arguments of the correct type.

If you want to use procedures defined in the book or reader as part of your solution to a programming problem, you must cite the page number on which it is defined so we know what you think it does.

## \*\*\* IMPORTANT \*\*\*

Check here if you are one of the people with whom we arranged to replace a missed/missing exam with other exam scores: \_\_\_\_\_

## \*\*\* IMPORTANT \*\*\*

If you have made grading complaints **that have not yet been resolved**, put the assignment name(s) here:

# **READ AND SIGN THIS:**

I certify that my answers to this exam are all my own work, and that I have not discussed the exam questions or answers with anyone prior to taking this exam.

If I am taking this exam early, I certify that I shall not discuss the exam questions or answers with anyone until after the scheduled exam time.

1-2	<u>/</u> 9
3–4	<u>/</u> 6
5	/4
6	/4
7	/7
8	/6
9	/6
10-11	/4
12	/6
13	/6
14	/0
15	/0
	/6
total	/70

## Question 1 (4 points):

(define (if-false x y) (if x 'true y))

For each of the following expressions, state the result in applicative order and in normal order. If the expression results in an error, just say ERROR; you don't have to give the exact message.

(if-false (/ 33 0) (/ 33 1))

Applicative:	Normal:
(if-false (/ 33 1) (/ 33 0))	

Applicative: \_\_\_\_\_

Normal: \_\_\_\_\_

### Question 2 (5 points):

Define a function maxnum that takes a list as argument. The elements of the list might or might not include numbers, along with other values. If any elements are numbers, return the largest number. If not, return **#f**. If an element is a list, don't look inside it; consider only top-level elements.

> (maxnum '(20 minutes and 45 seconds after 10 (2009 5 15)))
45

#### Use higher-order functions, not recursion!

Note: You get 4 points if your function works only for positive numbers.

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# Question 3 (4 points):

Here are two procedures that use two different algorithms to solve the same problem: Given two lists of length n as arguments, return true if they have any elements in common, false if not. For each procedure, indicate the order of growth in time, and whether the procedure generates an iterative or a recursive process.

```
(define (common? ls1 ls2)
  (cond ((null? ls1) #f)
        ((null? ls2) #f)
        ((equal? (car ls1) (car ls2)) #t)
        (else (or (common? ls1 (cdr ls2))
                   (common? (cdr ls1) ls2)))))
            \underline{\hspace{0.5cm}} \Theta(n) \qquad \underline{\hspace{0.5cm}} \Theta(n^2) \qquad \underline{\hspace{0.5cm}} \Theta(2^n)
\Theta(1)
 Iterative
                Recursive
(define (common? ls1 ls2)
  (cond ((null? ls1) #f)
        ((member (car ls1) ls2) #t)
        (else (common? (cdr ls1) ls2))))
            \Theta(1)
Iterative Recursive
Question 4 (2 points):
(define my-stream
  (cons-stream 3 (cons-stream 4 (add-streams my-stream
                                                  (stream-cdr my-stream)))))
```

```
> (show-stream my-stream 5)
```

## Question 5 (4 points):

Write a procedure sumpath that takes a Tree (with datum and children) of numbers as its argument, and returns a tree of the same shape in which each datum is the sum of all the data between the original datum and the root. For example, in the case below the value 5 in the original tree becomes 8 (1 + 2 + 5) in the returned tree:

## Question 6 (4 points):

(a) Fill in the blanks in the interactions below:

```
(define (swap ls1 ls2)
  (let ((temp ls1))
    (set! ls1 ls2)
    (set! ls2 temp) ) )
STk> (define foo (list 1 2 3))
foo
STk> (define bar (list 'a 'b 'c))
bar
STk> (swap foo bar)
okay
STk> foo
```

STk> bar

(b) Fill in the blank below:

STk> (define x 1) STk> (define y x) STk> (set! x 2) STk> y

### Question 7 (7 points):

We want to use OOP language (with define-class) to simulate a "smartphone" that combines a cell phone with a phonebook, calendar, etc. Assume that you are given a cellphone class, representing a bare-bones cell phone, that accepts a dial message with a phone number as argument.

```
> (define dumb-phone (instantiate cellphone))
> (ask dumb-phone 'dial '555-2368)
; ... phone dials
```

(a) Define a contact class that is instantiated with a person's name, address, and phone number, and accepts messages name, address, and phone that return the corresponding values:

```
> (define bh (instantiate contact 'Brian '(781 Soda) '642-8311))
> (ask bh 'phone)
642-8311
```

This question continues on the next page.

## Question 7 continued:

(b) Define a smartphone class that takes no instantiation arguments. The message addcontact takes three arguments for name, address, and phone number; it creates a contact object and adds it to its list of contacts. The message call takes a name as argument, finds the name in the list of contacts, and dials the corresponding phone number.

```
> (define my-phone (instantiate smartphone))
> (ask my-phone 'add-contact 'Dan '(777 Soda) '642-9595)
> (ask my-phone 'add-contact 'Brian '(781 Soda) '642-8311)
> (ask my-phone 'add-contact 'Mike '(779 Soda) '642-7017)
> (ask my-phone 'call 'Brian)
; ... phone dials 642-8311 (using dial method) ...
> (ask my-phone 'call 'Alonzo)
"Name not found"
```

### Question 8 (6 points):

What will the Scheme interpreter print in response to each of the following expressions? Also, <u>draw a "box and pointer" diagram</u> for the result of each printed expression. If any expression results in an error, just write "ERROR"; you don't have to give the precise message.

Hint: It'll be a lot easier if you draw the box and pointer diagram *first*!

```
(let ((x (list 'a 'b 'c)))
 (set-cdr! (cdr x) 3)
 x)
```

```
(let ((x (list 'a 'b 'c)))
 (set-car! x (caddr x))
 (set-car! (cddr x) (car x))
 x)
```

```
(let ((x (list 'a 'b 'c)))
 (set! (car x) (caddr x))
 x)
```

Question 9 (6 points):

```
(define x 100)
(define y 10)
(define s (make-serializer))
(define t (make-serializer))
(parallel-execute (s (lambda () (set! x (+ x y))))
                  (t (lambda () (set! y (* x y)))))
(list x y)
```

What are the possible correct answers for (list x y)?

What are the additional possible incorrect answers, if any?

Is deadlock possible? Yes No

```
(define x 100)
(define y 10)
(define s (make-serializer))
(define t (make-serializer))
(parallel-execute (s (t (lambda () (set! x (+ x y)))))
                  (s (t (lambda () (set! x (* x y)))))); x this time!
(list x y)
```

What are the possible correct answers for (list x y)?

What are the additional possible incorrect answers, if any?

Is deadlock possible? \_\_\_\_Yes No

#### Question 10 (2 points):

For each of the following interactions, indicate whether it will execute faster in the analyzing evalutor than in the original metacircular evalutor, or the same speed. Circle <u>Faster</u> or <u>Same</u> for each.

```
(define (mystery n)
  (if (< n 1)
      1
      (* n (mystery (/ n 2))))
> (mystery 100)
Analyzing will be:
                                    Same
                     Faster
(define (f x)
  (* (+ x 30) 7))
(define (g x)
  (* x x))
(define (h x)
  (/ (+ (* x 2) 10) 2))
(define (fgh n)
  (f (g (h n))))
> (fgh 100)
Analyzing will be:
                     Faster
                                    Same
```

#### Question 11 (2 points):

The following interaction is run in the *lazy* evaluator:

```
(define (darrentron a b c)
  (if a (* a b) c))
> (darrentron (* 1 2) (* 3 4) (* 1 2))
```

How many times is \* called if promises are memoized?

How many times is \* called if promises are *not* memoized?

# Question 12 (6 points):

This question concerns the *nondeterministic* (amb) evaluator.

(a) Write a procedure first-n that takes a list lst and a nonnegative integer num as argument, returning the first num elements of lst, or failing if there aren't enough elements.

(b) Using your answer to part (a), write a procedure sublist-of-length that takes a list and a nonnegative integer, and returns any sublist of num consecutive elements (not necessarily the first num elements). Each use of try-again should give another sublist, until there are no more of the required length.

### Question 13 (6 points):

Justin wants to improve his score in Rock Band, so he wants to find the song with the <u>lowest</u> average number of words per phrase. The file **rock-band-songs** has key-value pairs with a song title as the key and a phrase from the song as the value, like this:

```
(limelight . (living on a lighted stage))
(limelight . (approaches the unreal))
...
```

(a) Write a mapreduce call that will turn that file into a stream of key-value pairs, one per song, in which the value for each song is a list of two numbers: the total number of words in all the phrases of the song, and the number of phrases in the song. Call the returned stream countstream.

This question continues on the next page.

# Question 13 continued.

(b) Now write a mapreduce call that takes the stream from part (a) and returns a stream with song titles and average number of words per phrase for that song, sorted by the average, so that stream-car of the result will represent the song with the smallest average phrase length. Hint: The result from mapreduce is a stream that is sorted in key order; the smallest key comes first.

### Question 14 (6 points):

Write a set of rules for the query evaluator to implement the subst relation, with components old, new, input, and output. The first two are atoms; the last two can be atoms or (possibly deep) lists. The output should be the same as input except that every appearance of old, at any depth, is replaced by new:

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# Question 15 (6 points):

We want to add a new special form called let! to the metacircular evaluator. It takes a list of bindings, like let, but has no body. It doesn't make a new environment; instead, it changes the current environment. If there is already a binding for a variable named in the let! expression, then its value is changed, as in set!. If there isn't a binding, a new one is made in the current frame, as in define.

```
> (define count 1)
> (define (foo)
        (let! ((oldcount count) (count (+ count 1))))
        (* oldcount count))
> (foo)
2
> (foo)
6
> count
3
> oldcount
ERROR -- unbound variable: oldcount
```

The last expression is an error because oldcount was defined in foo's local environment, not in the global environment.

The relevant metacircular evaluator procedures are listed on the remaining pages of the exam. On this page, write the <u>names</u> of all the procedures that you modify elsewhere. New procedures can go here or on the back page.

```
(define (mc-eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((variable? exp) (lookup-variable-value exp env))
        ((quoted? exp) (text-of-quotation exp))
        ((assignment? exp) (eval-assignment exp env))
        ((definition? exp) (eval-definition exp env))
        ((if? exp) (eval-if exp env))
        ((lambda? exp)
         (make-procedure (lambda-parameters exp)
                         (lambda-body exp)
                         env))
        ((begin? exp)
         (eval-sequence (begin-actions exp) env))
        ((cond? exp) (mc-eval (cond->if exp) env))
        ((application? exp)
         (mc-apply (mc-eval (operator exp) env)
                   (list-of-values (operands exp) env)))
        (else
         (error "Unknown expression type -- EVAL" exp))))
(define (mc-apply procedure arguments)
  (cond ((primitive-procedure? procedure))
         (apply-primitive-procedure procedure arguments))
        ((compound-procedure? procedure)
         (eval-sequence
           (procedure-body procedure)
           (extend-environment
             (procedure-parameters procedure)
             arguments
             (procedure-environment procedure))))
        (else
         (error
          "Unknown procedure type -- APPLY" procedure))))
(define (definition? exp)
  (tagged-list? exp 'define))
(define (eval-definition exp env)
  (define-variable! (definition-variable exp)
                    (mc-eval (definition-value exp) env)
                    env)
 'ok)
```

```
(define (eval-assignment exp env)
 (set-variable-value! (assignment-variable exp)
                       (mc-eval (assignment-value exp) env)
                       env)
 'ok)
(define (make-frame variables values)
 (cons variables values))
(define (frame-variables frame) (car frame))
(define (frame-values frame) (cdr frame))
(define (add-binding-to-frame! var val frame)
  (set-car! frame (cons var (car frame)))
  (set-cdr! frame (cons val (cdr frame))))
(define (extend-environment vars vals base-env)
  (if (= (length vars) (length vals))
      (cons (make-frame vars vals) base-env)
      (if (< (length vars) (length vals))
          (error "Too many arguments supplied" vars vals)
          (error "Too few arguments supplied" vars vals))))
(define (lookup-variable-value var env)
 (define (env-loop env)
    (define (scan vars vals)
      (cond ((null? vars)
             (env-loop (enclosing-environment env)))
            ((eq? var (car vars))
             (car vals))
            (else (scan (cdr vars) (cdr vals)))))
    (if (eq? env the-empty-environment)
        (error "Unbound variable" var)
        (let ((frame (first-frame env)))
          (scan (frame-variables frame)
                (frame-values frame)))))
  (env-loop env))
```

```
(define (set-variable-value! var val env)
 (define (env-loop env)
    (define (scan vars vals)
      (cond ((null? vars))
             (env-loop (enclosing-environment env)))
            ((eq? var (car vars))
             (set-car! vals val))
            (else (scan (cdr vars) (cdr vals)))))
    (if (eq? env the-empty-environment)
        (error "Unbound variable -- SET!" var)
        (let ((frame (first-frame env)))
          (scan (frame-variables frame)
                (frame-values frame)))))
 (env-loop env))
(define (define-variable! var val env)
 (let ((frame (first-frame env)))
    (define (scan vars vals)
      (cond ((null? vars)
             (add-binding-to-frame! var val frame))
            ((eq? var (car vars))
             (set-car! vals val))
            (else (scan (cdr vars) (cdr vals)))))
    (scan (frame-variables frame)
          (frame-values frame))))
```