

## CS 172 , Spring, 1999 Professor M. Blum

### Problem #1

- Define the number of steps taken by a NDTM on input  $x$ .
- Define the number of steps taken by a NDTM on inputs of length of  $n$ .

a) # NDTMi  $[x] =$

case1:  $\min(y \text{ belongs to } (\text{summation } *)) \{ \# \text{DTMi}[y,x] \}$  ( this is the deterministic half of the NDTMi)

if there exists  $y$  such that  $\text{DTMi}[y,x]$  accepts (i.e. enters an accepting state)

case2: 1 otherwise

b) # NDTMi( $n$ ) =  $\text{Max}\{\# \text{NDTMi}[x]\}$ ,  $|x| = n$

### Problem #2

Define two (computational) problems  $p_1, p_2$  to be poly-time equivalent iff it is possible to solve  $p_1$  in polynomial time given an algorithm to solve  $p_2$  in polynomial time ( $p_1 \leq p_2$ ), and vice-versa ( $p_2 \leq p_1$ ).

Are the following two problems poly-time equivalent?

If so, prove it.

If not, explain why not.

Decision:

Instance: NDTMi,  $x$  in  $\{0,1\}^*$ ,  $m$  in unary

..... $m$

(ie  $1 \dots 1 = 1$  ).

Question: Does NDTMi accept  $x$  in  $m$  steps? ie does there exist a  $y$  in  $\{0,1\}^*$  s.t.  $\text{DTMi}$  accepts  $(y,x)$  in  $m$  steps, if any (ie if such  $y$  exists);

"NONE" if there is no such  $y$ .

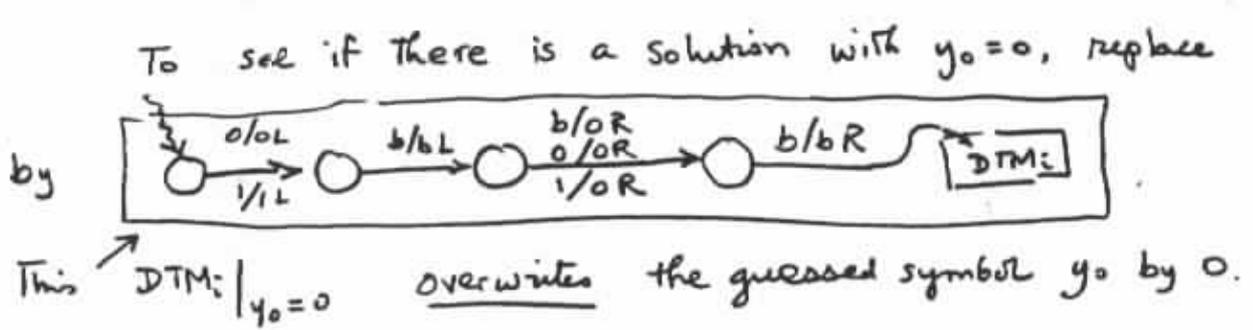
1) YES! Decision  $\leq$  Optimization:

- If optimization program returns  $y$ , then the Decision program returns YES
- If optimization program returns "NONE", then the DEcision program returns NO.

The running time of the Decision algorithm = running time of the Optimizatoin algorithm +  $O(1)$ .

2) Optimization  $\leq$  Decision:

- If Decision algor returns No, then the Opt. algor returns NONE.
- If Dec. algor returns YES, then the Opt. algor must find  $y$ . It can do this by finding the bits of  $y = Y_k Y_{(k-1)} \dots Y_0$  one at a time starting say with  $Y_0$ .



If the Dec. alg rejects  $[\text{NDTMil}(Y_0 = 0), x, m]$ , then we know that  $Y_0 = 1$ . Else we know that it's ok to let  $Y_0 = 0$ . In general, having determined  $Y(i-1) = A(i-1) \dots, Y_0 = A_0$ , one can determine  $Y_i$  by augmenting the  $\text{NDTM}_i$  to  $\text{NDTM}_{i+1} A(i-1) \dots A_0$ , which overwrites the guessed  $Y_i Y(i-1) \dots Y_0$  with  $A(i-1) \dots A_0$ . If the  $\text{NDTM}_{i+1} A(i-1) \dots A_0$  rejects, then we know that  $A_i = 1$ . Else ok to let  $A_i = 0$ . The size of the augmented  $\text{DTM}_{i+1} A(i-1) \dots A_0$  is just  $|\text{DTM}_i| + O(m)$ , which is poly in  $(|\text{NDTM}_i| + |x| + m)$ .

### Problem #3

Explain what problems if any you encounter in doing the above reductions in the case that  $m$  is given in binary instead of unary.

- Decision  $\leq$  Optimization as before, but
- Optimization (NOT  $\leq$ ) Decision:  
 The reason is that the required output  $y$  ( $y = 2^x$ ) can be terribly long, length  $|y| = x$ , for inputs of length  $n = |\text{NDTM}_i| + |x| + \lg m$  ( $m = \text{poly}(|y|)$ )  
 $= |\text{NDTM}_i| + |x| + O(|x|)$   
 $= \text{poly}(|x|)$ .  
 The decision algorithm can take just  $\text{poly}(n)$  steps, because it knows the existence of  $y$  without having to exhibit it. But the optimization algorithm must exhibit (print)  $y$ .

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**Posted by HKN (Electrical Engineering and Computer Science Honor Society)  
 University of California at Berkeley  
 If you have any questions about these online exams  
 please contact [examfile@hkn.eecs.berkeley.edu](mailto:examfile@hkn.eecs.berkeley.edu).**