

CS 172 , Spring, 1999 Professor M. Blum

Problem #1

- Define the number of steps taken by a NDTM on input x .
- Define the number of steps taken by a NDTM on inputs of length of n .

a) $\# \text{NDTM}_i[x] =$

case1: $\min(\text{y belongs to } (\text{summation } *)) \{ \# \text{DTM}_i[y,x] \}$ (this is the deterministic half of the NDTM_i)

if there exists y such that $\text{DTM}_i[y,x]$ accepts (i.e. enters an accepting state)

case2: 1 otherwise

b) $\# \text{NDTM}_i(n) = \text{Max}\{\# \text{NDTM}_i[x], |x| = n\}$

Problem #2

Define two (computational) problems p_1, p_2 to be poly-time equivalent iff it is possible to solve p_1 in polynomial time given an algorithm to solve p_2 in polynomial time ($p_1 \leq p_2$), and vice-versa ($p_2 \leq p_1$).

Are the following two problems poly-time equivalent?

If so, prove it.

If not, explain why not.

Decision:

Instance: $\text{NDTM}_i, x \text{ in } \{0,1\}^*, m \text{ in unary}$

..... m

(ie $1 \dots 1 = 1$).

Question: Does NDTM_i accept x in m steps? ie does there exist a y in $\{0,1\}^*$ s.t. DTM_i accepts (y,x) in m steps, if any (ie if such y exists);

"NONE" if there is no such y .

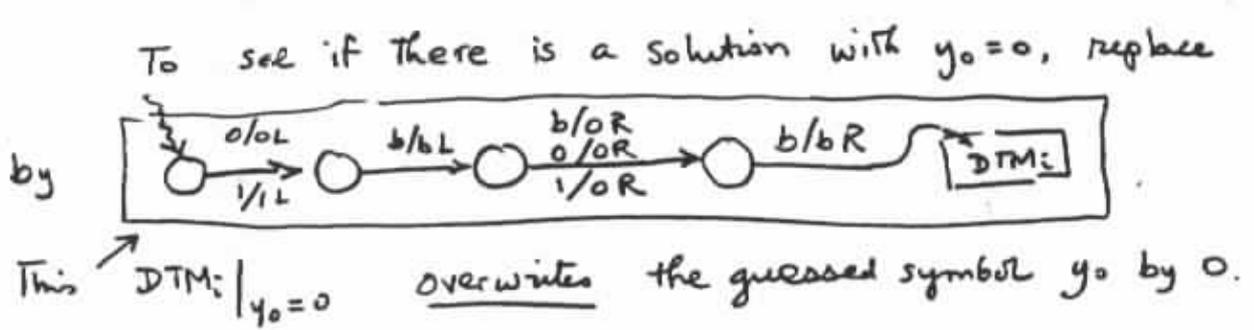
1) YES! Decision \leq Optimization:

- If optimization program returns y , then the Decision program returns YES
- If optimization program returns "NONE", then the DEcision program returns NO.

The running time of the Decision algorithm = running time of the Optimizaton algorithm + $O(1)$.

2) Optimization \leq Decision:

- If Decision algor returns No, then the Opt. algor returns NONE.
- If Dec. algor returns YES, then the Opt. algor must find y . It can do this by finding the bits of $y = Y_k Y_{(k-1)} \dots Y_0$ one at a time starting say with Y_0 .



If the Dec. alg rejects $[NDTMil(Y_0 = 0), x, m]$, then we know that $Y_0 = 1$. Else we know that it's ok to let $Y_0 = 0$. In general, having determined $Y(i-1) = A(i-1) \dots, Y_0 = A_0$, one can determine Y_i by augmenting the $NDTM_i$ to $NDTM_{i|A(i-1) \dots A_0}$, which overwrites the guessed $Y_i Y(i-1) \dots Y_0$ with $0A(i-1) \dots A_0$. If the $NDTM_{i|A(i-1) \dots A_0}$ rejects, then we know that $A_i = 1$. Else ok to let $A_i = 0$. The size of the augmented $DTM_{i|A_i \dots A_0}$ is just $|DTM_i| + O(m)$, which is poly in $(|NDTM_i| + |x| + m)$.

Problem #3

Explain what problems if any you encounter in doing the above reductions in the case that m is given in binary instead of unary.

- Decision \leq Optimization as before, but
- Optimization (NOT \leq) Decision:
 The reason is that the required output y ($y = 2^x$) can be terribly long, length $|y| = x$, for inputs of length $n = |NDTM_i| + |x| + \lg m$ ($m = \text{poly}(|y|)$)
 $= |NDTM_i| + |x| + O(|x|)$
 $= \text{poly}(|x|)$.
 The decision algorithm can take just $\text{poly}(n)$ steps, because it knows the existence of y without having to exhibit it. But the optimization algorithm must exhibit (print) y .

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