Midterm Exam #2
April 20, 2016
CS162 Operating Systems

Your Name: 

SID AND 162 Login: 

TA Name: 

Discussion Section Time: 

General Information:
This is a closed book and one 2-sided handwritten note examination. You have 80 minutes to answer as many questions as possible. The number in parentheses at the beginning of each question indicates the number of points for that question. You should read all of the questions before starting the exam, as some of the questions are substantially more time consuming.

Write all of your answers directly on this paper. Make your answers as concise as possible. If there is something in a question that you believe is open to interpretation, then please ask us about it!

Good Luck!!

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<tr>
<th>QUESTION</th>
<th>POINTS ASSIGNED</th>
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1. (12 points total) True/False and Why? CIRCLE YOUR ANSWER AND WRITE AN EXPLANATION (no credit will be given if no explanation is provided).

   a. In addition to greatly reducing the overall failure probability, moving from a single hard drive to using RAID 1 (mirroring) can also increase the throughput of small random reads.

   **TRUE**  
   **FALSE**  
   Why?

   b. The TCP transport protocol provides a reliable, in-order bytestream abstraction.

   **TRUE**  
   **FALSE**  
   Why?

   c. The Byzantine failure model allows arbitrary behavior by a faulty system component.

   **TRUE**  
   **FALSE**  
   Why?

   d. Remote Procedure Calls provide identical semantics to local calls.

   **TRUE**  
   **FALSE**  
   Why?
2. (26 points total) Memory and I/O.
   a. (12 points) For each statement, CIRCLE the address translation scheme(s) for which the statement is true.
      i) (4 points) A process cannot directly read or write another process’ memory. Assume there is no `mmap()` operation.
         No translation
         Base and Bound
         Segmentation
         Paging

      ii) (4 points) External fragmentation may occur.
          No translation
          Base and Bound
          Segmentation
          Paging

      iii) (4 points) Part of the process’ stack could be swapped out to disk while the process continues to execute using the rest of the stack. Assume the stack is located in a typical location for the type of address translation in use.
          No translation
          Base and Bound
          Segmentation
          Paging
b. (10 points) Consider a system with the following specifications:

- 46-bit virtual address space
- Page size of 8 KBytes
- Page table entry size of 4 Bytes
- Every page table is required to fit into a single page

How many levels of page tables would be required to map the entire virtual address space?

In the space below, document the format of a virtual address under this translation scheme. Briefly explain your rationale.

c. (4 points) Briefly explain, in four sentences or less, what is DMA and why DMA enables a system to more efficiently handle I/O devices.
3. (24 points total) Demand Paging.
   a. (18 points) Consider a demand paging system with 3 pages of physical memory. When a page fault occurs, we use a page replacement algorithm to make space for the new page by evicting a page from memory. Using the FIFO, Clock Algorithm, and LRU page replacement algorithms, mark which pages will be in physical memory after each page request. The first 3 rows for each algorithm are pre-filled for you.

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<tr>
<th>Time</th>
<th>FIFO</th>
<th>Clock</th>
<th>LRU</th>
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Number of page faults
b. (4 points) When a page fault occurs and a page is evicted, what entries in which data structures must the OS invalidate?

c. (2 points) In the Nth chance page replacement algorithm, explain the purpose of the parameter N. Consider the effects of a small or large N.
4. (22 points total) Disks and File Systems.
   a. (6 points) Name two performance-improving features that are provided by modern
      disk controllers. For each of the features you name, briefly explain the benefit
      consequences of a disk controller that does not provide that feature.
      i) Feature #1 – Benefits and consequences if it is not provided:

      ii) Feature #2 – Benefits and consequences if it is not provided:

   b. (5 points) If a disk subsystem receives an average of 10 requests per second and
      each request has an average response time of 500ms, how many requests are in
      the subsystem on average? Explain your answer.

   c. (4 points) The FAT file system eliminates external fragmentation by allocating disk
      space in block-sized units. Considering that fact, briefly explain why
      defragmentation is still sometimes necessary.
d. (3 points) In the UNIX 4.2 BSD FFS, inodes have direct pointers, a singly indirect pointer, a doubly indirect pointer, and a triply indirect pointer. The maximum file size supported by this inode type is approximately the same as the maximum file size supported by an inode with only a triply indirect pointer. Briefly explain one disadvantage of an inode design that only uses a triply indirect pointer instead of a combination of pointers (*other than the slightly reduced maximum file size*).

e. (4 points) Briefly, *in one to three sentences*, explain the purpose of the journal in a journaling file system?
5. (16 points total) Pintos coding question: Verifying user pointers.
Your project group is tasked with implementing a new syscall for Pintos. This syscall
takes a linked list as one of its arguments, and your responsibility is to create a
function that checks the validity of the linked list, so the syscall handler can safely
dereference it. The value of head may be NULL, if the list is empty. This is function
signature of your syscall:

\[
\text{void foo(struct node *head);}\
\]

We have also defined the structure of the node struct, as well as
\text{exit_thread_if_invalid()}, a function that will verify the validity of user-
specified pointers.

\text{struct node}
\{
    int number;
    char buffer[16];
    struct node *next;
    // If this is the last node, next = NULL
\}

/* Checks that P to P+SIZE-1 is a valid user buffer.
   Kills the current thread if it is invalid. */
\text{void exit_thread_if_invalid(void *p, size_t size)};

\text{uint32_t SYS_FOO = 60; // The syscall number for foo()}

Fill in the \text{syscall_handler()} function on the next page so that it \text{safely}
handles the SYS_FOO syscall.
static void syscall_handler (struct intr_frame *f) {
    uint32_t* args = ((uint32_t*) f->esp);

    // Check if this is SYS_FOO
    if (_______________________________) {
        // This will actually handle SYS_FOO,
        // assuming arguments are valid
        handle_foo(args, &f->eax);
    } else {
        // Code for ALL other syscalls will go here.
    }
}