

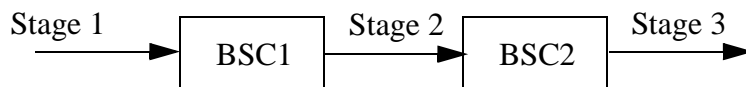
EECS 126 — Midterm #1

25 February 1999, 11:20 - 12:20 p.m.

Please explain clearly the steps in your answers.

Problem 1 (50 points)

Consider the cascade of two independent binary symmetric channels:



Let the error probabilities be ϵ_1 for BSC1 and ϵ_2 for BSC2. Suppose we generate a 0 or 1 randomly and input it to BSC1, observing the bit at all 3 stages.

Let B_i^x be the event that at stage i , an x is observed where $x \in \{0, 1\}$.

Assume $P(B_1^0) = P_0$ and that $P(B_1^x \cap B_3^y | B_2^z) = P(B_1^x | B_2^z) \cdot P(B_3^y | B_2^z)$ for $(x, y, z) \in \{0, 1\}^3$.

This says that B_1^x and B_3^y are conditionally independent given B_2^z .

Note that $P(B_2^0 | B_1^1) = P(B_2^1 | B_1^0) = \epsilon_1$ and $P(B_3^0 | B_2^1) = P(B_3^1 | B_2^0) = \epsilon_2$.

- Write down the sample space for this random experiment.
- Express the probability that we observe 0 at Stage 1, 1 at Stage 2, and 1 at Stage 3 in terms of the events $B_i^x, i \in \{1, 2, 3\}, x \in \{0, 1\}$.
- Show that $P(B_3^x | B_1^y \cap B_2^z) = P(B_3^x | B_2^z)$ using our conditional independence assumption.
- Calculate the probability in (b), above. (If you use the multiplication rule, (c) may be useful.)
- Calculate the probability that we observe a 0 at Stage 1 and a 1 at Stage 3.
- Calculate $P(B_3^1 | B_1^0)$.

Problem 2 (50 points)

Consider the following scheme for transmitting a data bit over a BSC.

1. Send the bit 3 times.
2. If 000 received, guess that 0 was transmitted.
If 111 received, guess that a 1 was transmitted.
Otherwise, request that the 3 bits be sent again.

Assume that errors occur independently every time the BSC is used, and that the error probability of the BSC is ϵ .

- a) What is the probability that the group of 3 bits must be sent k times before a decision is made?
- b) What is the probability that a decision is made after sending the group of 3 bits k times and it is the correct decision?
- c) What is the probability that the correct decision is made?
- d) Does this seem like an efficient scheme when $\epsilon = 0.1$ say?

(Hint: You might like to consider the expected number of bits that needs to be sent to send each bit of actual data.)

Can you think of a simple scheme using the expected number of bits sent per data bit that has a lower error probability?