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CS 70

Discrete Mathematics and Probability Theory

Summer 2014 James Cook

Midterm 1

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Thursday July 17, 2014, 12:40pm-2:00pm.

Instructions:

- Do not turn over this page until the proctor tells you to.
- Don't write any answers on the backs of pages (we won't be scanning those). There is an extra page at the end in case you run out of space.
- The exam has 9 pages (the last two are mostly blank).

PRINT your student ID: \_\_\_\_\_

PRINT AND SIGN your name: \_\_\_\_\_, \_\_\_\_\_, \_\_\_\_\_  
(last) (first) (signature)

PRINT your discussion section and GSI (the one you attend): \_\_\_\_\_

Name of the person to your left: \_\_\_\_\_

Name of the person to your right: \_\_\_\_\_

Name of someone in front of you: \_\_\_\_\_

Name of someone behind you: \_\_\_\_\_

PRINT your name and student ID: \_\_\_\_\_

## True/False

1. (16 pts.) For each of the following statements, circle T if it is true and F otherwise. You do not need to justify or explain your answers.

- T F For all positive integers  $x$  and  $p$ , if  $\gcd(x, p) = 1$ , then  $x^{p-1} \equiv 1 \pmod{p}$ .
- T F One way to prove a statement of the form  $P \implies Q$  is to assume  $\neg Q$  and prove  $\neg P$ .
- T F  $\forall x \exists y P(x, y) \equiv \exists x \forall y P(y, x)$ .
- T F  $P \implies (Q \implies R) \equiv (P \wedge Q) \implies R$
- T F  $P \implies (Q \wedge R) \equiv (P \implies Q) \vee R$
- T F To prove  $(\forall n \in \mathbf{N})P(n)$ , it is enough to prove  $P(0)$ ,  $P(2)$  and  $(\forall n \geq 2)(P(n) \implies P(n+2))$ .
- T F In a stable marriage instance, there can be two women with the same optimal man.
- T F In stable marriage, if Man 1 is at the top of Woman A's ranking but the bottom of every other woman's ranking, then every stable matching must pair 1 with A.

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## Short Answer

2. (4 pts.) Compute  $(2^3 \cdot 5^{71}) + (3^3 + 4^2) \bmod 8$ .

3. (4 pts.) Compute  $\frac{200 + 14 \cdot 102}{99} \bmod 10$ .

4. (4 pts.) Prove that  $(\exists x \in \mathbf{R})(\forall y \in \mathbf{R}) x \cdot y < 2$ .

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## RSA

5. (12 pts.) Someone sends Pandu an RSA-encrypted message  $x$ . The encrypted value is  $E(x) = 2$ . However, Pandu was silly and picked numbers far too small to make RSA secure. Given his public key  $(N = 77, e = 43)$ , find  $x$ .

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## Induction

6. (12 pts.) Prove that every two consecutive numbers in the Fibonacci sequence are coprime. (In other words, for all  $n \geq 1$ ,  $\gcd(F_n, F_{n+1}) = 1$ . Recall that the Fibonacci sequence is defined by  $F_1 = 1$ ,  $F_2 = 1$  and  $F_n = F_{n-2} + F_{n-1}$  for  $n > 2$ .)

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## Error-Correcting Codes

7. (15 pts.) Alice wants to send to Bob a message of length 3, and protect against up to 2 erasure errors. Using the error-correcting code we learned in class, she obtains a polynomial  $P(x)$  modulo 11 and sends 5 points to Bob. Bob only receives 3 of the points:  $P(1) = 4, P(3) = 1, P(4) = 5$ .
- (a) (12 pts.) Decode Alice's original message  $P(1), P(2), P(3)$ .

- (b) (3 pts.) If Alice tried to send a message with a modulus of 10 instead of 11, what exactly could go wrong? (You don't need to do any computations in your answer.)

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## Polynomials

8. (16 pts.) Suppose  $P$  is a polynomial over  $\mathbf{R}$ , and for every  $x, y \in \mathbf{R}$ ,  $P(x+y) = P(x) + P(y)$ .

(a) Prove that for every positive integer  $n$ ,  $P(n) = n \cdot P(1)$ .

(b) Prove that  $P$  has degree at most 1.

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[Extra page. If you want the work on this page to be graded, make sure you tell us on the problem's main page.]



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[Doodle page! Draw us something if you want or give us suggestions or complaints.]