

Midterm I - Solution

CS164, Spring 2014

March 3, 2014

- Please read all instructions (including these) carefully.
- **This is a closed-book exam. You are allowed a one-page handwritten cheat sheet.**
- Write your name, login, and SID.
- There are TODO pages in this exam and 3 questions, each with multiple parts. If you get stuck on a question move on and come back to it later.
- You have 1 hour and 15 minutes to work on the exam.
- Please write your answers in the space provided on the exam, and clearly mark your solutions. You may use the backs of the exam pages as scratch paper. **Do not** use any additional scratch paper.
- Solutions will be graded on correctness and *clarity*. Each problem has a relatively simple and straightforward solution. Partial solutions will be graded for partial credit.
- No electronic devices are allowed, including **cell phones** used merely as watches. Silence your cell phones and place them in your bag.

LOGIN: _____

NAME: _____

SID: _____

Problem	Max points	Points
1	28	
2	24	
3	48	
Sub Total	100	

1 Regular Expressions and Finite Automata

Consider a small language using only the letters “z”, “o”, and the slash character “/”. A comment in this language starts with “/o” and ends after the very next “o/”.

- (a) Give a regular expression that matches exactly one complete comment and nothing else. Assume comments do not nest. For full credit, use only the core notations: ϵ , “ab”, AB, A|B, and A^* **[8 points]**

$/o(o^*z|)^*o^*o/$

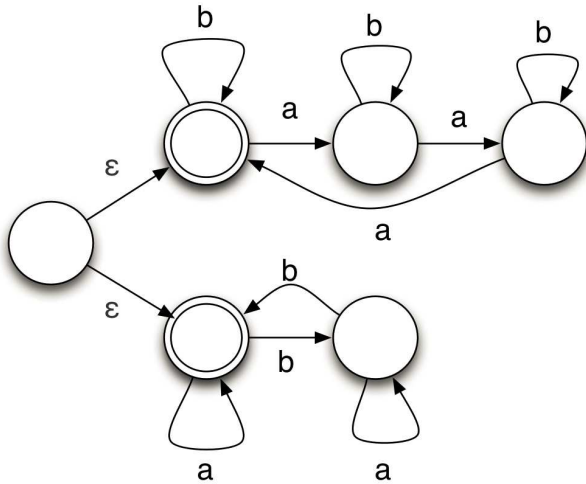
Consider the language over the alphabet $\Sigma = \{a,b\}$ containing strings in which number of 'a's is a multiple of 3

or number of 'b's is a multiple of 2.

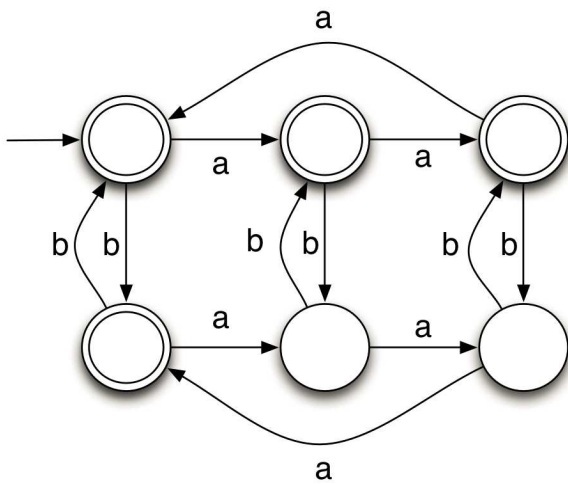
(b) Write a regular expression for this language. [4 points]

$(a^*ba^*ba^*)^* \mid (b^*ab^*ab^*ab^*)^*a^*b^*$

(c) Write an NFA for this language. [4 points]



(d) Write a DFA with at most 6 states for this language [4 points]



(3) Can we construct a regular expression for a language over the alphabet $\Sigma = \{a,b\}$ whose strings have equal number of occurrences of **a** and **b**? Explain. [4 points]

No. Corresponding state machine requires infinite number of states.

Yes/No (1)

Reason (3)

2 LL Parsing

Consider the following grammar with terminals $*$, $!$, n , $($, and $)$.

$$\begin{aligned} E &\rightarrow FH \\ H &\rightarrow *E \mid \epsilon \\ F &\rightarrow F! \mid G \\ G &\rightarrow n \mid (E) \end{aligned}$$

(a) The grammar is not LL(1). Explain in one sentence why [2 points]

$F \rightarrow F!$ is left recursive


(b) Fix the grammar to make it LL(1) by filling in the blanks below [4 points]

$$\begin{aligned} E &\rightarrow FH & F &\rightarrow \underline{\quad GK \quad} \\ H &\rightarrow *E \mid \epsilon & K &\rightarrow \underline{\quad !K!K \quad} \\ G &\rightarrow n \mid (E) & K &\rightarrow \underline{\quad \epsilon \quad} \end{aligned}$$

(c) Compute the first and the follow set of the fixed grammar [8 points]

	First	Follow
E	(, n	\$,)
F	(, n	\$,), *
G	(, n	\$,), *, !
H	*, ϵ	\$,)
K	!, ϵ	\$,), *

(d) Compute LL(1) parsing table [10 points]

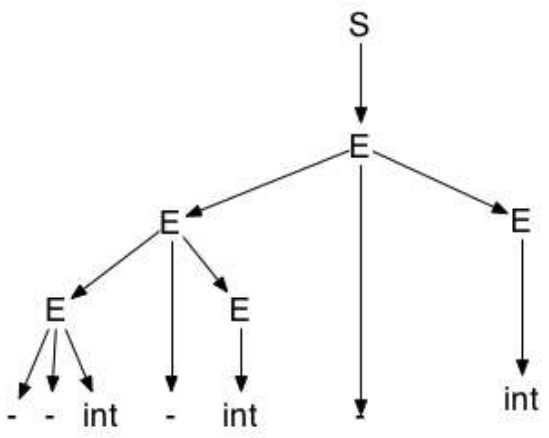
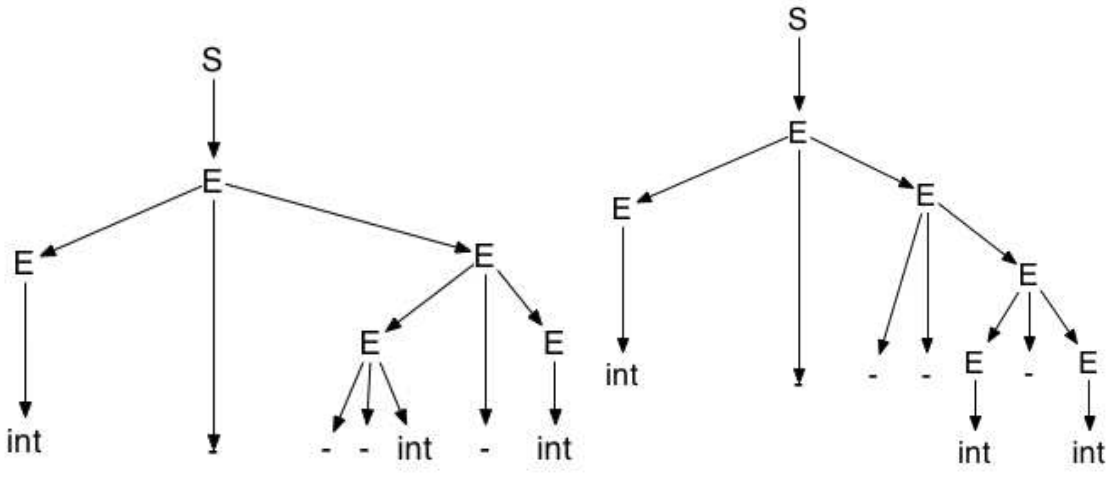
	*	!	()	\$	n
E			$E \rightarrow FH$			$E \rightarrow FH$
F			$F \rightarrow GK$			$F \rightarrow GK$
G			$G \rightarrow (E)$			$G \rightarrow n$
H	 $H \rightarrow *E$			ϵ	ϵ	
K	ϵ	$K \rightarrow !K$		ϵ	ϵ	

3 LR Parsing and Ambiguity

Consider the following grammar with terminals $-$ (the negation operator) and int .

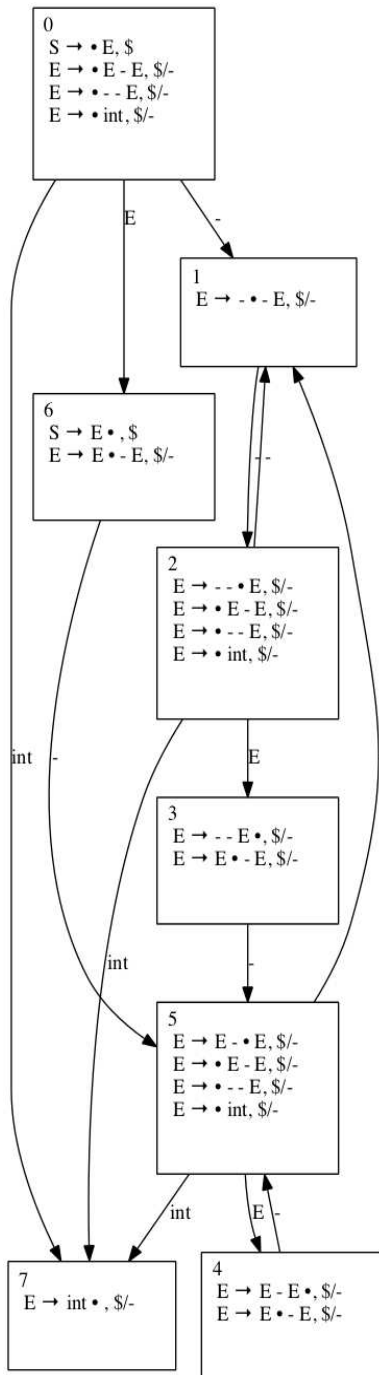
$$\begin{aligned} S &\rightarrow E \\ E &\rightarrow E - E \mid - - E \mid int \end{aligned}$$

(a) Draw all the parse trees for the string $int - - - int - int$ [6 points]



(b) Is this grammar ambiguous? Why or why not? [2 point]

The grammar is ambiguous because it admits multiple parse trees for string *int --- int - int*



(reduction labels)

- state 3: reduce $E \rightarrow -E$ on $\$/-$
- state 4: reduce $E \rightarrow E-E$ on $\$/-$
- state 6: accept $S \rightarrow E$ on $\$$
- state 7: reduce $E \rightarrow \text{int}$ on $\$/-$

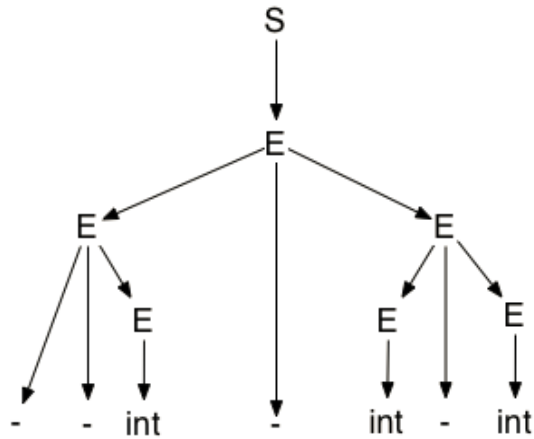
(c) Complete the above partial LR(1) DFA for the grammar. [16 points]

- Fill in items of all states by performing closure operation. (6)
- Fill in missing transition labels on all edges (4)
- Write the necessary “reduce by ... on ... ” labels on states (2)
- Add missing transition edges (*Hint: State 2 and State 5*) (4)

(d) For each state with a conflict, list the state, the lookahead token, and the type of conflict (i.e. shift-reduce conflict, or reduce-reduce conflict). [4 points]

3: Lookahead token: -, shift-reduce conflict
 4: Lookahead token: -, shift-reduce conflict

Suppose we want the string **- - int - int - int** to have only the following parse tree (call this property P).



(e) Describe in English the precedence and associativity rules necessary to ensure property P. [4 points]

1. “-” is right-associative
2. “--” has higher precedence than “-”

(f) Explain, for each conflict in the LR(1) parsing DFA for this grammar, how it should be resolved to ensure property P. [4 points]

State 3: reduce instead of shift on “-”
 State 4: shift instead of reduce on “-”

(g) Rewrite the grammar to an equivalent unambiguous grammar to ensure property P. Two grammars are equivalent when they accept the same language. [8 points]

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S → E
E → T - E | T
T → - - T | int
  
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(h) The Cool grammar has an ambiguity introduced by *let-expression*. Give an example illustrating the ambiguity associated with *let-expression*. [4 point]

let ... in 1 + 2

can be parsed as:

(let ... in 1) + 2

or

let ... in (1 + 2)